Computational Learning Theory: An Analysis of a Conjunction Learner

Machine Learning



Slides based on material from Dan Roth, Avrim Blum, Tom Mitchell and others

This lecture: Computational Learning Theory

- The Theory of Generalization
- Probably Approximately Correct (PAC) learning
- Positive and negative learnability results
- Agnostic Learning
- Shattering and the VC dimension

Where are we?

- The Theory of Generalization
 - When can be trust the learning algorithm?
 - What functions can be learned?
 - Batch Learning
- Probably Approximately Correct (PAC) learning
- Positive and negative learnability results
- Agnostic Learning
- Shattering and the VC dimension

This section

- 1. Analyze a simple algorithm for learning conjunctions
- 2. Define the PAC model of learning
- 3. Make formal connections to the principle of Occam's razor

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- 2. Define the PAC model of learning
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The true function $f = x_2 \land x_3 \land x_4 \land x_5 \land x_{100}$

Training data

- <(1,1,1,1,1,1,...,1,1), 1>
- <(1,1,1,0,0,0,...,0,0), 0>
- <(1,1,1,1,0,...0,1,1), 1>
- <(1,0,1,1,1,0,...0,1,1), 0>
- <(1,1,1,1,0,...0,0,1), 1>
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 $f = x_2 \wedge x_3 \wedge x_4 \wedge x_5 \wedge x_{100}$

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A simple learning algorithm (*Elimination*)

• Discard all negative examples

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A simple learning algorithm (*Elimination*)

- Discard all negative examples
- Build a conjunction using the features that are common to all positive conjunctions

$$h = x_1 \wedge x_2 \wedge x_3 \wedge x_4 \wedge x_5 \wedge x_{100}$$

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- <(1,1,1,1,1,1,...,1,1), 1>
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- <(1,1,1,1,1,1,....**0**1), 1>

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Positive examples *eliminate* irrelevant features

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A simple learning algorithm:

- Discard all negative examples
- Build a conjunction using the features that are common to all positive conjunctions

$$h = x_1 \land x_2 \land x_3 \land x_4 \land x_5 \land x_{100}$$

Clearly this algorithm produces a conjunction that is consistent with the data, that is $err_s(h) = 0$, if the target function is a monotone conjunction Exercise: Why?

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Claim 1: Any hypothesis consistent with the training data will only make mistakes on positive future examples Why?

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A mistake will occur only if some literal z (in our example x_1) is present in h but not in f

This mistake can cause a positive example to be predicted as negative by h Specifically: $x_1 = 0$, $x_2 = 1$, $x_3 = 1$, $x_4 = 1$, $x_5 = 1$, $x_{100} = 1$

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The reverse situation can never happen For an example to be predicted as positive in the training set, every relevant literal must have been present

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Theorem: Suppose we are learning a conjunctive concept with n dimensional Boolean features using m training examples. If

$$m > \frac{n}{\epsilon} \left(\log(n) + \log\left(\frac{1}{\delta}\right) \right)$$

then, with probability > 1 - δ , the error of the learned hypothesis err_D(h) will be less than ϵ .

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Poly in n, 1/ δ , 1/ ϵ

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If we see these many training examples, then the algorithm will produce a conjunction that, with high probability, will make few errors

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Let's prove this assertion

 $f = x_2 \wedge x_3 \wedge x_4 \wedge x_5 \wedge x_{100} \qquad \qquad h = x_1 \wedge x_2 \wedge x_3 \wedge x_4 \wedge x_5 \wedge x_{100}$

What kinds of examples would drive a hypothesis to make a mistake?

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What kinds of examples would drive a hypothesis to make a mistake?

Positive examples, where x_1 is absent f would say true and h would say false

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Let's quantify our surprise at seeing such examples

Let p(z) be the probability that, in an example drawn from D, the feature z is absent but the example has a positive label

- That is, after training is done, p(z) is the probability that in a randomly drawn example, the literal z causes a mistake
- For any z in the target function, p(z) = 0

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p(x₁): Probability that this situation occurs

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We know that
$$err_D(h) \leq \sum_{z \in h} p(z)$$

Via direct application of the union bound

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Union bound

For a set of events, probability that at least one of them happens < the sum of the probabilities of the individual events

n = dimensionality

- Call a literal z bad if $p(z) > \frac{\epsilon}{n}$
- Intuitively, a bad literal is one that has a significant probability of not appearing with a positive example
 - (And, if it appears in all positive training examples, it can cause errors)

If there are no bad literals, then $err_D(h) < \epsilon$

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- Why? Because
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Let us try to see when this will not happen

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What if there are bad literals?

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What if there are bad literals?

- Let z be a bad literal
- What is the probability that it will not be eliminated by one training example?

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Pr(z survives one example) = 1 - Pr(z is eliminated by one example)

$$\leq 1 - p(z) \ < 1 - rac{\epsilon}{n}$$

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There was one example of this kind

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What we know so far:

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There are at most n bad literals. So

 $Pr(\text{Any bad literal survives } m \text{ examples}) < n \left(1 - \frac{\epsilon}{n}\right)^m$

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We want this probability to be small

Why? So that we can choose enough training examples so that the probability that any z survives all of them is less than some δ

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That is, we want
$$n\left(1-\frac{\epsilon}{n}\right)^m < \delta$$

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We know that $1 - x < e^{-x}$. So it is sufficient to require $ne^{-\frac{m\epsilon}{n}} < \delta$

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Or equivalently,
$$m > \frac{n}{\epsilon} \left(\log(n) + \log\left(\frac{1}{\delta}\right) \right)$$

To guarantee a probability of failure (i.e, error > ϵ) that is less than δ , the number of examples we need is

$$m > \frac{n}{\epsilon} \left(\log(n) + \log\left(\frac{1}{\delta}\right) \right)$$

Poly in n, 1/
$$\delta$$
, 1/ ϵ

That is, if m has this property, then

- With probability 1 δ , there will be no bad literals,
- Or equivalently, with probability 1 δ , we will have err_D(h) < ϵ

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 Poly in n, 2

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What does this mean:

• If $\epsilon = 0.1$ and $\delta = 0.1$, then for n = 100, we need 6908 training examples

 $1/\delta$, $1/\epsilon$

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What does this mean:

- If ϵ = 0.1 and δ = 0.1, then for n = 100, we need 6908 training examples
- If ϵ = 0.1 and δ = 0.1, then for n = 10, we need only 461 examples

To guarantee a probability of failure (i.e, error > ϵ) that is less than δ , the number of examples we need is

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- If ϵ = 0.1 and δ = 0.1, then for n = 10, we need only 461 examples
- If ϵ = 0.1 and δ = 0.01, then for n = 10, we need 691 examples

To guarantee a probability of failure (i.e, error > ϵ) that is less than δ , the number of examples we need is

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What we have here is a PAC guarantee

Our algorithm is Probably Approximately Correct